

# Notes on BCFW Recursion in MATHEMATICA

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ABSTRACT: Basic reference formulary and key functions know for coding BCFW recursion in MATHEMATICA. Prepared for the *Amplitudes 2018* Summer School.

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## 1 Symbolic BCFW Recursion in Momentum-Twistor Space

We will denote the  $n$ -point  $N^k$ MHV tree-amplitude as ‘ $\mathcal{A}_n^{(k)}$ ’. When expressed in terms of momentum-twistor variables, MHV ( $k=0$ ) amplitudes are the identity for all multiplicity,  $\mathcal{A}_n^{(0)}=1$ , and amplitudes vanish when  $k < 0$  or  $k > (n-4)$  (for  $n > 3$ ).

The momenta of the particles involved in an amplitude are encoded by momentum (super)twistors denoted ‘ $z_a$ ’. When it is important or necessary to make clear *which* set of momentum twistors are involved in an amplitude, this will be indicated by writing ‘ $\mathcal{A}_n^{(k)}(a, b, \dots)$ ’ for  $\mathcal{A}_n^{(k)}(z_a, z_b, \dots)$ .

The BCFW recursion relations allow us to determine the  $n$ -point  $N^k$ MHV tree-amplitude  $\mathcal{A}_n^{(k)}$  in terms of lower-point trees as follows:

$$\begin{aligned} \mathcal{A}_n^{(k)}(1, \dots, n) &= \mathcal{A}_{n-1}^{(k)}(1, \dots, n-1) \\ &+ \sum_{\substack{n_L+n_R=n+2 \\ k_L+k_R=k-1}} \mathcal{A}_{n_L}^{(k_L)}(1, \dots, a-1, \hat{a}) R[1, a-1, a, n-1, n] \mathcal{A}_{n_R}^{(k_R)}(\hat{a}, a, \dots, n-1, \hat{n}), \end{aligned} \quad (1.1)$$

where the ‘shifted’ twistors  $\hat{a}$  and  $\hat{n}$  are defined according to:

$$\begin{aligned} \hat{a} &\equiv (a \ a-1) \cap (n-1 \ n \ 1) \equiv z_a \langle a-1 \ n-1 \ n \ 1 \rangle + z_{a-1} \langle n-1 \ n \ 1 \ a \rangle; \\ \hat{n} &\equiv (n \ n-1) \cap (a-1 \ a \ 1) \equiv z_n \langle n-1 \ a-1 \ a \ 1 \rangle + z_{n-1} \langle a-1 \ a \ 1 \ n \rangle. \end{aligned} \quad (1.2)$$

Because the recursion relations successively lower  $k$  (or  $n$ ) and all MHV ( $k=0$ ) amplitudes are 1, the only non-trivial functions that survive successive recursion via (1.1) are the so-called ‘ $R$ -invariants’ defined according to:

$$R[a, b, c, d, e] \equiv \frac{\delta^{1 \times 4} (\eta_a \langle bcde \rangle + \eta_b \langle cdea \rangle + \eta_c \langle deab \rangle + \eta_d \langle eabc \rangle + \eta_e \langle abcd \rangle)}{\langle a \ b \ c \ d \rangle \langle b \ c \ d \ e \rangle \langle c \ d \ e \ a \rangle \langle d \ e \ a \ b \rangle \langle e \ a \ b \ c \rangle}, \quad (1.3)$$

where  $\{z_a, \dots, z_e\}$  can be related to external momentum twistors through (possibly nested) shifts according to (1.2).

## 2 How to Do Supersymmetry (without Anticommutation)

Superfunctions multiply in the obvious way; and so a product of  $k$   $R$ -invariants will take the form of  $5k$  4-brackets  $\langle \dots \rangle$  in the denominator with  $k$   $(1 \times 4)$  fermionic  $\delta$ -functions in the numerator. Fermionic  $\delta$ -functions also simply multiply, but because<sup>1</sup>

$$\int d\eta_a d\eta_b \delta(a_1\eta_a + b_1\eta_b)\delta(a_2\eta_a + b_2\eta_b) = (a_1b_2 - a_2b_1), \quad (2.1)$$

for anti-commuting variables, it is most natural to consider the  $k \times n$  matrix ‘ $C$ ’ of  $\eta$ -coefficients (with one row of  $C$  coming from each  $R$ -invariant) and think of these superfunctions to be in the form,

$$f(z) \delta^{k \times 4}(C \cdot \eta) \equiv f(z) \prod_{I=1}^4 \prod_{\alpha=1}^k \left( \sum_{a=1}^n C_a^\alpha \eta_a^I \right), \quad (2.2)$$

where  $f(z)$  is an ordinary (bosonic) function, and  $C$  is a  $k \times n$  matrix of ordinary functions. As illustrated in the simple case in (2.1), the reason for thinking of superfunctions in this way is that all component functions—obtained by projecting out (via integration)  $4k$  of the  $\eta$ ’s will result in simply multiplication of  $f(z)$  by  $4$   $(k \times k)$  determinants of the matrix  $C$ . To see this, notice that for each of the  $SU(4)$ -indices  $I \in \{1, \dots, 4\}$ ,  $\eta$ -integration results in a determinant,

$$\int d\eta_{a_1}^I \cdots d\eta_{a_k}^I \prod_{\alpha=1}^k \left( \sum_{a=1}^n C_a^\alpha \eta_a^I \right) = (a_1 \cdots a_k) \equiv \det(C_{a_1}^\alpha, \dots, C_{a_k}^\alpha). \quad (2.3)$$

(This is a natural generalization of (2.1).) Thus, we see that the product of  $k$   $R$ -invariants corresponds to a superfunction with  $\binom{n}{k}^4$  bosonic components—corresponding to a choice of a  $(k \times k)$  minor of the matrix  $C$  for each  $I \in \{1, \dots, 4\}$ .

The point of this aside is that we will consider superfunctions (in MATHEMATICA) to correspond to a pair of purely bosonic data,

$$f(z) \delta^{k \times 4}(C(z) \cdot \eta) \Leftrightarrow \{f(z), C(z)\}, \quad (2.4)$$

from which all component functions can be extracted by multiplying  $f(z)$  by  $4$   $(k \times k)$  minors of the matrix  $C(z)$ . For superfunctions that are products of  $R$ -invariants,  $f(z)$  is the product of their bosonic parts, and  $C(z)$  is a matrix whose  $k$  rows are generated one-by-one from each  $R$ -invariant in the product.

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<sup>1</sup>Recall that fermionic  $\delta$ -functions are trivial:  $\int d\eta \delta(f(\eta)) \equiv \int d\eta f(\eta)$  for any Grassmann  $\eta$ .

### 3 Glossary of MATHEMATICA Functions to Know (and Love)

The following is a list of MATHEMATICA functions used in the live-demonstration of the implementation of tree-level BCFW (listed roughly in order of appearance).

- **Replace**[*x*,*y*] (same as '*x/.y*') and **ReplaceRepeated**[*x*,*y*] (same as '*x//.y*')
- **Rule**[*x*,*y*] (same as '*x->y*') and **RuleDelayed**[*x*,*y*] (same as '*x:>y*')
- **Set**[*x*,*y*] (same as '*x=y*') and **SetDelayed**[*x*,*y*] (same as '*x:=y*')
- **Which**[*l*] (see also **If**[*l*])
- **And**[*x*,*y*] (same as '*x&& y*') and **Or**[*x*,*y*] (same as '*x|| y*')
- **Equal**[*x*,*y*] (same as '*x==y*') and **SameQ**[*x*,*y*] (same as '*x===y*'), etc.
- **With**[*l*], (see also **Block**[*l*], **Module**[*l*])
- **Table**[**function**[*index*], {*index*,*indexRange*}] (see also **Sum**[*l*], **Prod**[*l*])
- **Join**[*l*] (see also **Flatten**[*l*])
- **Select**[*list*, **testFunction**[*l*]]: selects those elements of *list* that are mapped to **True** under **testFunction**.
- **Apply**[**function**[*l*], *exprn*] (same as '**function**@@*exprn*'—see also '**f**@@*exprn*')  
e.g. **f**@@{*x*,*y*} returns **f**[*x*,*y*]; **f**@@@{{*x*,*y*},{*z*,*w*}} returns {**f**[*x*,*y*], **f**[*z*,*w*]}
- **Map**[**function**[*l*], *list*] (same as: '**function**/@*list*')
- **Function**[{*x*}, **function**[*x*]] (same as '**function**[#]&')
- **Total**[*l*], **Times**[*l*], **Plus**[*l*], etc.
- **ReplacePart**[*l*]
- **Range**[*l*], **Partition**[*l*], **RotateLeft**[*l*], etc.
- **SparseArray**[*l*] (and **Normal**[*l*])
- **Thread**[*l*] (especially in the context of **Thread**[**Rule**[*l*]])
- **Cases**[*l*]
- **Det**[*l*] (please remember to never reinvent **Det**[*l*]!)
- **RandomInteger**[*l*]
- ...

If any of the functions above are unfamiliar, consult MATHEMATICA's *Documentation Center* for more details (and illustrative examples). This can be accessed quickly by prepending a '?' to a function in the notebook interface (e.g. by typing '**?Partition**'). You will often find that many built-in functions have optional additional arguments (that can be extremely useful to know!).

Finally, if you think that '**Function**' is a bit odd or irrelevant, you are wrong: (especially when constructed using the '**(#)&@**' syntax) **Function** is one of the most useful programming constructs in MATHEMATICA—and the same goes for **Map** (**/@**) and **Apply** (**@@** and **@@@**). Get to know and love their syntax and usage—it will dramatically improve your programming prowess.